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NOTIFICATION OF ELECTION

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Date of mailing (day/month/year) 16 August 2000 (16.08.00)	Applicant's or agent's file reference GIC-562 PCT
International application No. PCT/US99/25973	Priority date (day/month/year) 04 December 1998 (04.12.98)
International filing date (day/month/year) 04 November 1999 (04.11.99)	
Applicant LING, Fan	

1. The designated Office is hereby notified of its election made:

☒ in the demand filed with the International Preliminary Examining Authority on:
06 June 2000 (06.06.00)

☐ in a notice effecting later election filed with the International Bureau on:

2. The election ☒ was

☐ was not

made before the expiration of 19 months from the priority date or, where Rule 32 applies, within the time limit under Rule 32.2(b).

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INTERNATIONAL SEARCH REPORT

(PCT Article 18 and Rules 43 and 44)

Applicant's or agent's file reference GIC-562 PCT	FOR FURTHER ACTION see Notification of Transmittal of International Search Report (Form PCT/ISA/220) as well as, where applicable, item 5 below.	
International application No. PCT/US 99/ 25973	International filing date (day/month/year) 04/11/1999	(Earliest) Priority Date (day/month/year) 04/12/1998
Applicant GENERAL INSTRUMENT CORPORATION et al.		

This International Search Report has been prepared by this International Searching Authority and is transmitted to the applicant according to Article 18. A copy is being transmitted to the International Bureau.

This International Search Report consists of a total of 3 sheets.

☒ It is also accompanied by a copy of each prior art document cited in this report.

1. Basis of the report

- a. With regard to the language, the international search was carried out on the basis of the international application in the language in which it was filed, unless otherwise indicated under this item.

☐ the international search was carried out on the basis of a translation of the international application furnished to this Authority (Rule 23.1(b)).

- b. With regard to any nucleotide and/or amino acid sequence disclosed in the international application, the international search was carried out on the basis of the sequence listing:

☐ contained in the international application in written form.

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☐ the statement that the subsequently furnished written sequence listing does not go beyond the disclosure in the international application as filed has been furnished.

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2. ☐ Certain claims were found unsearchable (See Box I).

3. ☐ Unity of invention is lacking (see Box II).

4. With regard to the title,

☒ the text is approved as submitted by the applicant.

☐ the text has been established by this Authority to read as follows:

5. With regard to the abstract,

☒ the text is approved as submitted by the applicant.

☐ the text has been established, according to Rule 38.2(b), by this Authority as it appears in Box III. The applicant may, within one month from the date of mailing of this international search report, submit comments to this Authority.

6. The figure of the drawings to be published with the abstract is Figure No.

☒ as suggested by the applicant.

☐ because the applicant failed to suggest a figure.

☐ because this figure better characterizes the invention.

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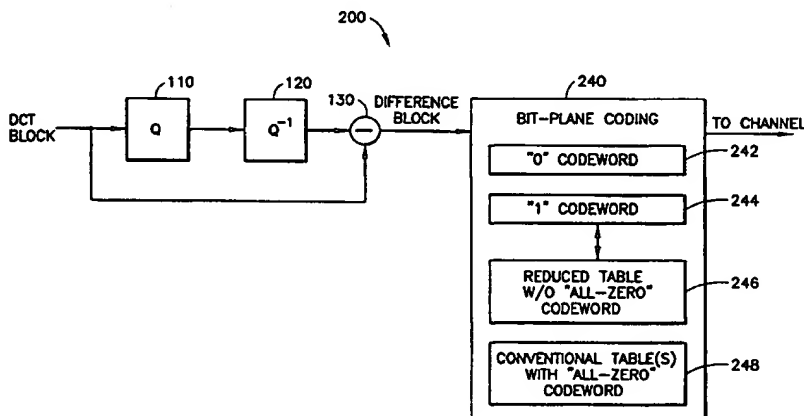
☐ None of the figures.



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(54) Title: IMPROVEMENT OF FINE GRANULARITY SCALABILITY USING BIT PLANE CODING OF TRANSFORM COEFFICIENTS

**(57) Abstract**

A system for efficient bit plane coding of transform coefficient data, such as DCT data used in a video coding system. Decimal values for the transform coefficients, e.g., in a block of several coefficients, are converted to binary values, where each bit occupies a corresponding bit plane, from the most significant bit to the least significant bit. One bit from each coefficient is provided in a common bit plane. A one-bit flag or codeword (such as "0") is used for coding - one or more initial all-zero bit planes, while another one-bit flag (such as "1") is used for designating the first subsequent non-all-zero plane. For the first non-all-zero plane, a reduced coding table is used to provide codewords that follow the one-bit flag. The coding table is reduced in size since it does not require a special "all-zero" codeword. Additionally, the use of a one-bit flag for designating the initial all-zero bit planes reduces the required number of coding bits over prior art schemes that require multi-bit all-zero codewords. An encoder (200) includes a "0" codeword function (242), a "1" codeword function (244), a reduced table (246), and conventional tables (248). A corresponding decoder (400) includes a "0" codeword function (442), a "1" codeword function (444), a reduced table (446), and conventional tables (448).

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**IMPROVEMENT OF FINE GRANULARITY SCALABILITY USING BIT
PLANE CODING OF TRANSFORM COEFFICIENTS**

BACKGROUND OF THE INVENTION

5 This application claims the benefit of U.S.
Provisional Application No. 60/110,882, filed December
4, 1998.

10 The present invention relates to a method and
apparatus for efficient bit plane coding of transform
coefficients, such as Discrete Cosine Transform (DCT)
coefficients. Such coefficients may be used in a
variety of applications, including digital video
encoding and decoding. In particular, an improvement
upon the technique known as Fine Granularity
Scalability using Bit plane (FGSB) coding is presented.

15 FGSB coding keeps the base layer coding technique
untouched. For example, the base layer coding
technique can be MPEG-2, MPEG-4, or any DCT-based
image/video coding technique. In the base layer, the
DCT coefficients are coded using a relatively coarse
20 quantization to obtain low bit rate data.

25 FIG. 1 illustrates a prior art apparatus for Fine
Granularity Scalability using Bit plane (FGSB) coding.
With FGSB coding, a difference (or residue) is obtained
between the original integer DCT coefficients and
dequantized DCT coefficients. As shown in the encoder
100, for example, an original block of DCT coefficients
is quantized at a quantizer 110, then the quantized

coefficients are dequantized (i.e., inverse quantized) at an inverse quantizer 120 to obtain the dequantized DCT coefficients. A difference block of DCT coefficients is output from a subtracter 130 and
5 provided to a bit plane coding function 140 before being communicated across a channel. For example, the data may be communicated in a broadband communication network, such as a cable or satellite television network, or a computer network, such as a local area
10 network (LAN), metropolitan area network (MAN), wide area network (WAN), internet, intranet, and the Internet. As mentioned, the base layer DCT coefficient data is coded conventionally.

The bit plane coding function 140 includes one or
15 more tables 145 for coding the bits in each bit plane. The tables include an "all-zero" codeword that is used when a bit plane has all zeroes values.

Since each DCT difference block typically has a few bit planes (e.g., up to four to eight-planes in
20 typical applications), fine granularity is achieved at very low complexity cost. The number of bit planes is determined by the number of bits needed to code the largest difference value. Essentially, DCT difference data in the successive bit plane layers can be used to
25 reduce the DCT coefficient quantization error. One or more of the bit plane layers can be recovered by a decoder according to the available channel bandwidth and the decoder's processing speed. The bit plane layers are recovered starting from the layer carrying
30 the most significant bit of the DCT difference data,

then the layer carrying the next most significant bit of the DCT difference data, and so forth.

FGSB coding can be simplified into the following steps:

- 5 1. After performing base layer coding, which is
DCT based, take the difference between the original DCT
coefficients and the dequantized DCT coefficients.
Find the number of bit planes required to code this
difference block.
- 10 2. Find out the maximum number of bit planes from
all the difference blocks of one video frame.
3. Code the maximum number of bit planes at the
very beginning of one frame's enhancement layer
bitstream.
- 15 4. Sequentially code the bit planes of one frame
starting from the Most Significant Bit (MSB) level.
5. When coding a bit plane, 2-D symbols are
formed of two components. The first component
indicates the number of consecutive zeroes (e.g., zero
20 run-length) until the next "1.". The second component
is a one-bit flag that indicates if there are any "1"s
left in the current bit plane. The second component is
therefore an End-Of-Plane (EOP) indicator. These 2-D
symbols are then entropy coded. If there are no "1"s
25 at all in the current bit plane, an "all-zero" symbol
is coded.

The following illustrates an example of coding of a particular bit plane using the above approach, designated "prior art 1".

Position:	0	1	2	3	4	5	6	...	63
bit value:	1	0	0	1	0	1	0	0's	0
prior art 1, 2-D symbol	(0,0)	(2,0)			(1,1)				

In the first row, the "position" is the sequence of the bit in the bit plane. For example, for an 8x8 bit plane, there are 64 bits, e.g., 0-63. In the second row, the bit value is shown, either a binary 0 or 1.

In the third row, the 2-D symbols used by the prior art scheme discussed above are shown. Specifically, at position 0, the bit value is "1", and the symbol is (0,0). The first component of the symbol, "0", indicates that the run-length of zero symbols is 0 (i.e., there is no zero) before the next "1". The second component of the symbol, "0", which is the EOP indicator, indicates that there is at least one subsequent "1" in the bit plane. Note that the symbols are given in decimal numbers, which are later converted to binary (e.g., $2_{10}=10_2$).

At position 1, the bit value is "0", and the symbol is (2,0). The first component of the symbol, "2", indicates that the run-length of zero symbols is 2 (i.e., there are zeroes in positions 1 and 2). Thus, the distance to the next "1" is 2 bit values (i.e., position 3). The second component of the symbol, "0",

again indicates that there are additional "1's" in the bit plane (i.e., in addition to the "1" in position 3).

At position 4, the bit value is "0", and the symbol is (1,1). The first component of the symbol, "1", indicates that the run-length of zero symbols is 1 (i.e., there is a zero in position 4). Thus, the distance to the next "1" is 1 bit value (i.e., position 5). The second component of the symbol, "1", indicates that there are no additional "1's" in the bit plane (after position 5).

The above FGS coding method is discussed in ISO/IEC JTC1/SC29/WG11, MPEG98/M4204, December 1998, "Fine Granularity Scalability Using Bit plane Coding of DCT Coefficients." With this technique, a bit plane with all zeroes is coded using the "all-zero" symbol, regardless of which bit plane layer is being coded. The all-zero symbol may need to be coded for more than one bit plane layer for a given DCT block. This is disadvantageous due to the size of the all-zero symbol, and the fact that it increases the coding table size.

In particular, the size (e.g., bit length) of the all-zero symbol is dictated by the entropy coding (e.g., Huffman coding) performed on the bit plane's 2-D symbols. As is known, when coding source symbols that are not equally probable, it is efficient to use variable-length code words. The probability of occurrence of the source symbols is used to select the code words so that more probable source symbols are assigned a shorter code word.

With these constraints, the length of the all-zero symbol is typically two bits or more. Furthermore,

since there are several thousand 8x8 DCT blocks in an image (e.g., consider a 525x480 pixel NTSC image), the data overhead caused by the all-zero symbol is significant.

5 Accordingly, it would be desirable to provide a method and apparatus for efficient bit plane coding which improves over the above techniques. The system should reduce the number of bits required to indicate the presence of a bit plane in which all values are
10 zero. The system should avoid the need for multiple "all-zero" symbols in the initial MSB bit plane layers of a single block having all zeroes, thereby reducing the data overhead for coding the bit plane.

15 The system should improve coding efficiency by reducing the number of symbols in a MSB level entropy coding table, thereby reducing the code length of the remaining symbols in the table.

20 The system should be compatible with coding schemes that provide multiple coding tables which are used for the different bit plane layers. Since the probability of occurrence of specific bits is different for the different bit plane layers, coding efficiency can be optimized by tailoring the coding table to the bit plane layer. This concept is discussed further in
25 the Ph.D. dissertation entitled "Optimization of Entropy Coding Efficiency Under Complexity Constraints in Image and Video Compression," (1998), Section 4.3, by Dr. Fan Ling, catalogued at the Electrical Engineering Department, Lehigh University,
30 Pennsylvania, USA.

The present invention provides a system having the above and other advantages.

SUMMARY OF THE INVENTION

The present invention relates to a method and apparatus for efficient bit plane coding of transform coefficients.

5 A method for efficient coding of a plurality of bit planes in which transform coefficient data is carried, includes the step of providing a codeword having a first state (e.g., "0") designating that a bit plane has all binary zeroes, and a second state (e.g.,
10 "1") designating that a bit plane does not have all binary zeroes. The most significant bit (MSB) bit plane is coded with the "0" when the MSB bit plane has all binary zeroes. Proceeding from the MSB bit plane toward the least significant bit (LSB) bit plane, each
15 successive bit plane that has all binary zeroes is coded with the "0", if any such successive bit planes are present, until a first bit plane that does not have all binary zeroes is reached. The first bit plane is coded with the "1", followed by at least one codeword
20 obtained from a first entropy coding table according to the bits in the first bit plane.

The codeword is preferably a one-bit codeword.

Importantly, the first entropy coding table does not include a multi-bit codeword for coding a bit plane
25 with all binary zeroes. Thus, the size of this coding table can be reduced relative to prior art schemes, which, in turn, reduces the number of coding bits.

A second, conventional entropy coding table is provided for coding one or more bit planes that follow

the first bit plane. This coding table includes a multi-bit codeword for coding a bit plane with all binary zeroes, and is therefore not reduced in size.

5 The transform coefficient data may include Discrete Cosine Transform (DCT) data, and/or image data.

10 A corresponding decoding method includes the step of (a) providing a decoding function for a codeword having a first state (e.g., "0") designating that a bit plane has all binary zeroes, and a second state (e.g., "1") designating that a bit plane does not have all binary zeroes. Proceeding from the MSB bit plane toward the least significant bit (LSB) bit plane, the "0" is decoded for each successive bit plane that has
15 all binary zeroes, if any such successive bit planes are present, until a first bit plane that does not have all binary zeroes is reached. The "1" is decoded for the first bit plane, and a first entropy decoding table is used for decoding at least one codeword that follows
20 the "1".

A related digital signal, and encoder and decoder apparatuses are also presented.

BRIEF DESCRIPTION OF THE DRAWINGS

FIG. 1 illustrates a prior art apparatus for Fine Granularity Scalability coding using Bit plane (FGSB) coding.

5 FIG. 2 illustrates an apparatus for Fine Granularity Scalability coding using Bit plane (FGSB) coding in accordance with the present invention.

FIG. 3 illustrates a coding method in accordance with the present invention.

10 FIG. 4 illustrates a decoder in accordance with the present invention.

DETAILED DESCRIPTION OF THE INVENTION

The present invention relates to a method and apparatus for efficient bit plane coding of transform coefficients.

5 The present invention is illustrated with reference to the table below, in which the bit plane coding symbols discussed in ISO/IEC JTC1/SC29/WG11, MPEG98/M4204, discussed above, are designated by "prior art 2", and bit plane coding symbols of the present
10 invention are designated by "invention".

 An example 8x8 DCT difference block as follows is assumed:

Position	:	0	1	2	3	4	5	6	7	. . .	63
Decimal Value:		11	1	10	3	2	1	0	0	. . .	0

15 Bit planes:

MSB-0	:	0	0	0	0	0	0	. . .	0	(all zeroes)
MSB-1	:	1	0	1	0	0	0	. . .	0	
MSB-2	:	0	0	0	0	0	0	. . .	0	(all zeroes)
MSB-3	:	1	0	1	1	1	0	. . .	0	
20 MSB-4	:	1	1	0	1	0	1	. . .	0	

"Prior art 2" Approach:

MSB-0 : (all-zero symbol),
MSB-1 : (0,0), (1,1),
MSB-2 : (all-zero symbol),
5 MSB-3 : (0,0), (1,0), (0,0), (0,1)
MSB-4 : (0,0), (0,0), (1,0), (1,1)

Present invention:

MSB-0 : 0 (one-bit flag=true),
MSB-1 : 1 (one-bit flag=false), (0,0), (1,1),
10 MSB-2 : (all-zero),
MSB-3 : (0,0), (1,0), (0,0), (0,1)
MSB-4 : (0,0), (0,0), (1,0), (1,1)

The base or MSB bit plane layer bits are designated "MSB-0," the second MSB bit plane layer bits are designated "MSB-1," the third MSB bit plane layer bits are designated "MSB-2," the fourth MSB bit plane layer bits are designated "MSB-3," and the fifth MSB bit plane layer bits are designated "MSB-4." MSB-4 is essentially the Least Significant Bit (LSB) bit plane layer.
20

The decimal values of the DCT difference coefficients are provided in positions 0-63. The corresponding binary values are provided in the different layers. For example, the decimal value 11 corresponds to the binary value 01011,, where the first "0" (reading from left to right) is the MSB. This "0" is therefore provided in MSB-0. The remaining binary digits are distributed in MSB-1 ("1"), MSB-2 ("0"), MSB-3 ("1"), and MSB-4 ("1").
25

Note that all the MSB-layer 0 values in this example are zeroes, indicating that the decimal values of the coefficients in the current block are all less than fifteen. The five bit plane layers allow decimal values up to 31. Of course, fewer or more bit planes may be used.

Since different DCT blocks have different numbers of bit planes, the MSB layer bit planes have a high probability of being all zero. The present invention takes advantage of this condition.

Coding using the prior art 2 technique in the above example is first discussed. MSB-0 is coded using the "all-zero" symbol, which typically requires two or more coded bits. This indicates the layer is all zeroes. For example, a coding table may be used that assigns a code word to a symbol. The number of coded bits will vary based on the coding table size and the entropy coding scheme used.

MSB-1 is coded using the (0,0) symbol, where the first "0" indicates a zero run length of zero (i.e., no zeroes) before the next "1", and the second "0" indicates there are additional "1's" left in the bit plane after the next "1". Next, the (1,1) symbol is used, where the first "1" indicates a zero run length of one (i.e., one zero) before the next "1", and the second "1" indicates there are no additional "1's" left in the bit plane after the next "1".

MSB-2 is coded using the "all-zero" symbol, indicating the layer is all zeroes.

MSB-layer 3 ("1") is coded using the (0,0) symbol. Next, position 1 ("0") is coded using the (1,0) symbol,

position 3 ("1") is coded using the (0,0) symbol, and position 4 ("1") is coded using the (0,1) symbol.

Position 0 of MSB-4 ("1") is coded using the (0,0) symbol, position 1 ("1") is coded using the (0,0) symbol, position 2 ("0") is coded using the (1,0) symbol, and position 4 ("0") is coded using the (1,1) symbol.

Accordingly, the prior art 2 technique uses the "all-zero" symbol twice, in MSB-0 and MSB-2. This technique is not optimal since the "all-zero" symbol has a length of two bits or more, and the presence of this symbol increases the length of the coding table. In contrast, the present invention allows the use of a reduced coding table, which does not have the all-zero symbol as an entry, for coding of MSB-0 and MSB-1.

Specifically, with the present invention, a one-bit flag is introduced to indicate if the MSB bit plane is all zeroes or not. For example, "0" can indicate all zeroes, and a "1" can indicate not all zeroes, or vice-versa. However, whenever a "1" is coded, the flag is not coded for all the following lower level bit planes of the block. Thus, the coding scheme of the invention is applied to the initial MSB layer (layer 0) when it has all zeroes, to any immediately following layers that have all zeroes, if such layers are present, and to the first following layer that does not have all zeroes.

If any further all-zero bit planes are encountered (e.g., after the bit plane encoded with the "1" flag), they are coded using the "all-zero" symbol of prior art 2.

This, in MSB-0, the (0) codeword or flag is used to indicate that the MSB-layer 0 is all zeroes. In MSB-1, (1)(0,0) designates that the codeword "1" is to be followed in the bitstream by the codeword for (0,0), and the codeword for (1,1), which latter two codewords are obtained from the reduced coding table.

The total amount of bits is reduced relative to the prior art 2 scheme since a reduced size coding table is used that does not have the all-zero symbol as an entry. This bit savings can be significant when multiplied over the many transform blocks. For example, consider the number of 8x8 DCT blocks in a 720x480 television picture.

FIG. 2 illustrates an apparatus for Fine Granularity Scalability coding using Bit plane (FGSB) coding in accordance with the present invention.

Like-numbered elements correspond to one another in the figures.

Here, a revised bit plane coding function 240 is used, which includes a "0" codeword function 242 for coding a "0" codeword for an all-zero bit plane as discussed herein. In the above example, MSB-0 is coded by this function 242.

A "1" codeword function 244 is provided for coding a "1" codeword for a first non-zero bit plane as discussed herein.

A reduced table 246 does not include an "all-zero" codeword and therefore has a reduced length relative to a conventional table. The reduced length results in a reduced number of coded bits for at least some of the table entries. The reduced table is associated with

the "1" codeword function 244 since it is used for coding bits in the bit plane that are associated with the "1" codeword. For example, in the above example, MSB-1 is coded by the functions 244 and 246.

5 Specifically, the symbols (0,0) and (1,1) in MSB-1 are coded by the function 246.

A conventional table function 248 includes one or more tables with the "all-zero" codeword. These tables therefore have an increased length relative to the
10 reduced table 246, and will use additional bits to code a bit plane. One or more conventional tables may be used. If multiple conventional tables are used, each table may be tailored to specific bit planes, as discussed in the Background. Three conventional tables
15 is believed to be suitable for FGSB video applications. Generally, the probability distribution of bits in the lower bit plane layers become more similar for each additional layer, so the advantage of a separate conventional coding table is reduced.

20 FIG. 3 illustrates a coding method in accordance with the present invention. At block 305, the first (i.e., most significant) bit plane of a new DCT block is processed. If the initial layer, MSB layer-0 has all zeroes (block 310), the layer is coded with a
25 codeword "0" (block 315) in accordance with the invention.

If the initial layer does not have all zeroes, processing continues at block 325.

30 Following block 315, if there are any remaining layers (block 318), a determination is made as to whether any immediately following layers also have all

zeroes (block 320), in which case they are also coded using the "0" symbol (block 315) in accordance with the invention.

5 If there are no remaining layers (block 318), the next DCT block is processed (block 305).

At block 325, if the next layer does not have all zeroes, a codeword "1" is provided, followed by one or more codewords for symbol pairs (e.g., 0,0, 0,1, 1,0 or 1,1) from a reduced table. In accordance with the
10 invention, the reduced table does not have an all-zero symbol, so the length of the code words for at least some of the symbols is reduced relative to the conventional table. Thus, the total amount of bits needed to code a transform block is reduced.

15 If there are any remaining layers (block 330), the next layer is processed (block 335) using the conventional table (block 340).

FIG. 4 illustrates a decoder in accordance with the present invention. The decoder 400 includes a bit
20 plane decoding function 440, which receives a coded data stream from a channel. For example, the decoder 400 may be a set-top terminal in a cable or satellite television network for receiving digital video data, e.g., according to the MPEG standard.

25 The "0" codeword (decoding) function 442 is a counterpart of the coding function 242 at the encoder 200, and identifies a coded zero as a designator that the bit plane contains all zeroes. The "1" codeword (decoding) function 444 is a counterpart of the coding
30 function 244 at the encoder 200, and identifies a coded

one as a designator of a first non-zero plane following a plane coded with a "0" codeword.

5 The reduced table 446 is a decoding function that is a counterpart of the table 246 at the encoder 200, and identifies symbol pairs that correspond to codewords in the received data stream for the plane that is also identified by a coded "1".

10 The conventional table(s) 448 are decoding functions that are counterparts of the tables 248 at the encoder 200, and identify symbol pairs that correspond to codewords in the received data stream for the plane(s) that are not identified by a coded "1".

15 The bit plane decoding function 440 outputs transform difference block data for further, conventional processing, the details of which will be apparent to those skilled in the art.

20 Note that the invention is even more efficient for coding smaller decimal values, where there are several initial MSB layers with all zeroes. Generally, at least one coded bit may be saved for each layer that is coded in accordance with the invention. Typically, most of the DCT difference values are zero or close to zero, with only a few larger values whose symbols require the use of the initial MSB layers. However, in
25 any case, coding efficiency is improved since the number of symbols from the MSB level entropy coding table is reduced by one symbol. As a result, the code length of the remaining symbols is also reduced.

30 This can be illustrated by a simplified example, by assuming a symbol set with three symbols, {A, B, C}.

Their occurring probabilities and Huffman codes are as follows:

	symbol:	A	B	C
	Probability:	0.4	0.3	0.3
5	Huffman code:	1	01	00
	code length (bits):	1	2	2

If we take one symbol (e.g., A) out of the set, the new set will be {B,C}. The above table becomes a reduced table:

10	symbol:	B	C
	Probability:	0.5	0.5
	Huffman code:	0	1
	code length (bits):	1	1

15 Comparing the two cases, we find that by taking one symbol out of the symbol set, we reduce the code length of all the remaining symbols. This conclusion applies generally to any entropy coding technique, and is not limited to Huffman coding. Note that the overall per cent savings in coded bits will be less as
20 the number of symbols, and the symbol length, in a coding table increases. However, even a small savings for each transform block is significant when multiplied by the many blocks in an image, for example.

25 Accordingly, it can be seen that the present invention provides a method and apparatus for efficient bit plane coding that provides a one-bit flag or codeword for coding one or more initial all-zero bit

planes and the first subsequent non-all-zero plane.
For the first non-all-zero plane, a reduced entropy
coding table is used to provide codewords that follow
the one-bit flag. The coding table is reduced in size
5 since it does not require a special "all-zero"
codeword. Additionally, the use of a one-bit flag for
designating the initial all-zero bit planes reduces the
required number of coding bits over prior art schemes
that require multi-bit all-zero codewords. The
10 invention thereby reduces the data overhead for coding
the bit plane.

Although the invention has been described in
connection with various specific embodiments, those
skilled in the art will appreciate that numerous
15 adaptations and modifications may be made thereto
without departing from the spirit and scope of the
invention as set forth in the claims.

The invention is suitable for use with any bit
plane coding scheme, and is not limited to bit plane
20 coding of DCT coefficients. For example, the invention
may be used with other transform coding techniques,
such as the Discrete Fourier Transform, Karhunen-Loeve
Transform, Walsh Hadamard Transform, wavelet transform,
as well as other known spatial transforms.

25 Moreover, any type of entropy coding, such as
Huffman coding or arithmetic coding, which uses a
floating point code length, may be used.

Moreover, the invention is not limited to coding
of MPEG-coded data or video data, but may be used,
30 e.g., for coding seismic, vibrational, temperature,

pressure and other types of data with 2-D or higher characteristics.

Additionally, the invention may be implemented using any known hardware, firmware, and/or software techniques.

What is claimed is:

1. A method for efficient coding of a plurality of bit planes in which transform coefficient data is carried, comprising the steps of:

(a) providing a codeword having a first state designating that a bit plane comprises all binary zeroes, and a second state designating that a bit plane does not comprise all binary zeroes;

(b) coding the most significant bit (MSB) bit plane with said codeword in its first state when the MSB bit plane comprises all binary zeroes;

(c) proceeding from the MSB bit plane toward the least significant bit (LSB) bit plane, coding each successive bit plane that comprises all binary zeroes with said codeword in its first state, if any such successive bit planes are present, until a first bit plane that does not comprise all binary zeroes is reached; and

(d) coding said first bit plane with said codeword in its second state, followed by at least one codeword obtained from a first entropy coding table according to the bits in said first bit plane.

2. The method of claim 1, wherein:
the codeword is a one-bit codeword.

3. The method of claim 1, wherein:
the first entropy coding table does not include a multi-bit codeword for coding a bit plane with all binary zeroes.

4. The method of claim 1, comprising the further step of:

providing a second entropy coding table for coding at least one bit plane that follows the first bit plane; wherein:

the second entropy coding table includes a multi-bit codeword for coding a bit plane with all binary zeroes.

5. The method of claim 1, wherein:

the transform coefficient data comprises Discrete Cosine Transform (DCT) data.

6. The method of claim 1, wherein:

the transform coefficient data comprises image data.

7. A method for decoding a plurality of bit planes in which transform coefficient data is carried, comprising the steps of:

(a) providing a decoding function for a codeword having a first state designating that a bit plane comprises all binary zeroes, and a second state designating that a bit plane does not comprise all binary zeroes;

wherein the most significant bit (MSB) bit plane is coded with said codeword in its first state when the MSB bit plane comprises all binary zeroes;

(b) proceeding from the MSB bit plane toward the least significant bit (LSB) bit plane, decoding said

codeword in its first state for each successive bit plane that comprises all binary zeroes, if any such successive bit planes are present, until a first bit plane that does not comprise all binary zeroes is reached; and

(c) decoding said codeword in its second state for said first bit plane, then using a first entropy decoding table for decoding at least one codeword that follows said codeword in its second state;

wherein the at least one codeword is obtained from a first entropy coding table according to the bits in said first bit plane.

8. The method of claim 7, wherein:
the codeword is a one-bit codeword.

9. The method of claim 7, wherein:
the first entropy decoding table does not include a multi-bit codeword for decoding a bit plane with all binary zeroes.

10. The method of claim 7, comprising the further step of:

providing a second entropy decoding table for decoding at least one bit plane that follows the first bit plane; wherein:

the second entropy decoding table includes a multi-bit codeword for decoding a bit plane with all binary zeroes.

11. The method of claim 7, wherein:
the transform coefficient data comprises Discrete Cosine Transform (DCT) data.

12. The method of claim 7, wherein:
the transform coefficient data comprises image data.

13. A digital signal carrying data for efficiently coding of a plurality of bit planes in which transform coefficient data is carried, comprising:

(a) a codeword having a first state designating that a bit plane comprises all binary zeroes, and a second state designating that a bit plane does not comprise all binary zeroes; wherein:

the most significant bit (MSB) bit plane is coded with said codeword in its first state when the MSB bit plane comprises all binary zeroes; and

proceeding from the MSB bit plane toward the least significant bit (LSB) bit plane, each successive bit plane that comprises all binary zeroes is coded with said codeword in its first state, if any such successive bit planes are present, until a first bit plane that does not comprise all binary zeroes is reached;

(b) said codeword in its second state for coding said first bit plane; and

(c) at least one codeword following said codeword in its second state obtained from a first entropy

coding table according to the bits in said first bit plane.

14. The signal of claim 13, wherein:
the codeword is a one-bit codeword.

15. An apparatus for efficient coding of a plurality of bit planes in which transform coefficient data is carried, comprising:

(a) means for providing a codeword having a first state designating that a bit plane comprises all binary zeroes, and a second state designating that a bit plane does not comprise all binary zeroes;

(b) means for coding the most significant bit (MSB) bit plane with said codeword in its first state when the MSB bit plane comprises all binary zeroes;

(c) means for coding each successive bit plane that comprises all binary zeroes with said codeword in its first state, if any such successive bit planes are present, proceeding from the MSB bit plane toward the least significant bit (LSB) bit plane, until a first bit plane that does not comprise all binary zeroes is reached; and

(d) means for coding said first bit plane with said codeword in its second state, followed by at least one codeword obtained from a first entropy coding table according to the bits in said first bit plane.

16. The apparatus of claim 15, wherein:
the codeword is a one-bit codeword.

17. The apparatus of claim 15, wherein:
the first entropy coding table does not include a multi-bit codeword for coding a bit plane with all binary zeroes.

18. The apparatus of claim 15, further comprising:

a second entropy coding table for coding at least one bit plane that follows the first bit plane;
wherein:

the second entropy coding table includes a multi-bit codeword for coding a bit plane with all binary zeroes.

19. An apparatus for decoding a plurality of bit planes in which transform coefficient data is carried, comprising:

(a) a decoding function for a codeword having a first state designating that a bit plane comprises all binary zeroes, and a second state designating that a bit plane does not comprise all binary zeroes;

wherein the most significant bit (MSB) bit plane is coded with said codeword in its first state when the MSB bit plane comprises all binary zeroes;

(b) means for decoding said codeword in its first state for each successive bit plane that comprises all binary zeroes, if any such successive bit planes are present, proceeding from the MSB bit plane toward the least significant bit (LSB) bit plane, until a first bit plane that does not comprise all binary zeroes is reached; and

(c) means for decoding said codeword in its second state for said first bit plane, then using a first entropy decoding table for decoding at least one codeword that follows said codeword in its second state;

wherein the at least one codeword is obtained from a first entropy coding table according to the bits in said first bit plane.

20. The apparatus of claim 19, wherein:
the codeword is a one-bit codeword.

21. The apparatus of claim 19, wherein:
the first entropy decoding table does not include a multi-bit codeword for decoding a bit plane with all binary zeroes.

22. The apparatus of claim 19, further comprising:

a second entropy decoding table for decoding at least one bit plane that follows the first bit plane;
wherein:

the second entropy decoding table includes a multi-bit codeword for decoding a bit plane with all binary zeroes.

AMENDED CLAIMS

[received by the International Bureau on 08 June 2000 (08.06.00);
Original claims 2, 8, 14, 16 and 20 amended; remaining claims unchanged (4 pages)]

1. A method for efficient coding of a plurality of bit planes in which transform coefficient data is carried, comprising the steps of:

(a) providing a codeword having a first state designating that a bit plane comprises all binary zeroes, and a second state designating that a bit plane does not comprise all binary zeroes;

(b) coding the most significant bit (MSB) bit plane with said codeword in its first state when the MSB bit plane comprises all binary zeroes;

(c) proceeding from the MSB bit plane toward the least significant bit (LSB) bit plane, coding each successive bit plane that comprises all binary zeroes with said codeword in its first state, if any such successive bit planes are present, until a first bit plane that does not comprise all binary zeroes is reached; and

(d) coding said first bit plane with said codeword in its second state, followed by at least one codeword obtained from a first entropy coding table according to the bits in said first bit plane.

2. The method of claim 1, wherein:

the codeword having the first and second states is a one-bit codeword.

3. The method of claim 1, wherein:

the first entropy coding table does not include a multi-bit codeword for coding a bit plane with all binary zeroes.

codeword in its first state for each successive bit plane that comprises all binary zeroes, if any such successive bit planes are present, until a first bit plane that does not comprise all binary zeroes is reached; and

(c) decoding said codeword in its second state for said first bit plane, then using a first entropy decoding table for decoding at least one codeword that follows said codeword in its second state;

wherein the at least one codeword is obtained from a first entropy coding table according to the bits in said first bit plane.

8. The method of claim 7, wherein:

the codeword having the first and second states is a one-bit codeword.

9. The method of claim 7, wherein:

the first entropy decoding table does not include a multi-bit codeword for decoding a bit plane with all binary zeroes.

10. The method of claim 7, comprising the further step of:

providing a second entropy decoding table for decoding at least one bit plane that follows the first bit plane; wherein:

the second entropy decoding table includes a multi-bit codeword for decoding a bit plane with all binary zeroes.

coding table according to the bits in said first bit plane.

14. The signal of claim 13, wherein:
the codeword having the first and second states is a one-bit codeword.

15. An apparatus for efficient coding of a plurality of bit planes in which transform coefficient data is carried, comprising:

(a) means for providing a codeword having a first state designating that a bit plane comprises all binary zeroes, and a second state designating that a bit plane does not comprise all binary zeroes;

(b) means for coding the most significant bit (MSB) bit plane with said codeword in its first state when the MSB bit plane comprises all binary zeroes;

(c) means for coding each successive bit plane that comprises all binary zeroes with said codeword in its first state, if any such successive bit planes are present, proceeding from the MSB bit plane toward the least significant bit (LSB) bit plane, until a first bit plane that does not comprise all binary zeroes is reached; and

(d) means for coding said first bit plane with said codeword in its second state, followed by at least one codeword obtained from a first entropy coding table according to the bits in said first bit plane.

16. The apparatus of claim 15, wherein:
the codeword having the first and second states is a one-bit codeword.

(c) means for decoding said codeword in its second state for said first bit plane, then using a first entropy decoding table for decoding at least one codeword that follows said codeword in its second state;

wherein the at least one codeword is obtained from a first entropy coding table according to the bits in said first bit plane.

20. The apparatus of claim 19, wherein:
the codeword having the first and second states is a one-bit codeword.

21. The apparatus of claim 19, wherein:
the first entropy decoding table does not include a multi-bit codeword for decoding a bit plane with all binary zeroes.

22. The apparatus of claim 19, further comprising:
a second entropy decoding table for decoding at least one bit plane that follows the first bit plane;
wherein:
the second entropy decoding table includes a multi-bit codeword for decoding a bit plane with all binary zeroes.

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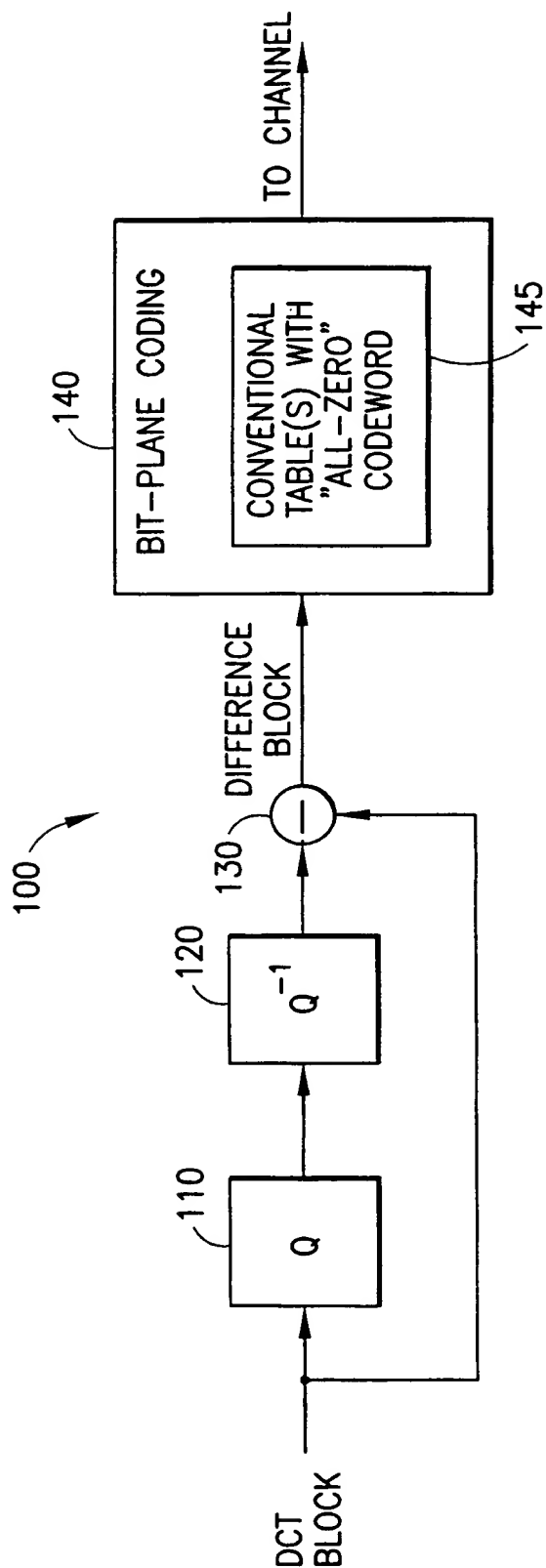


FIG.1
PRIOR ART

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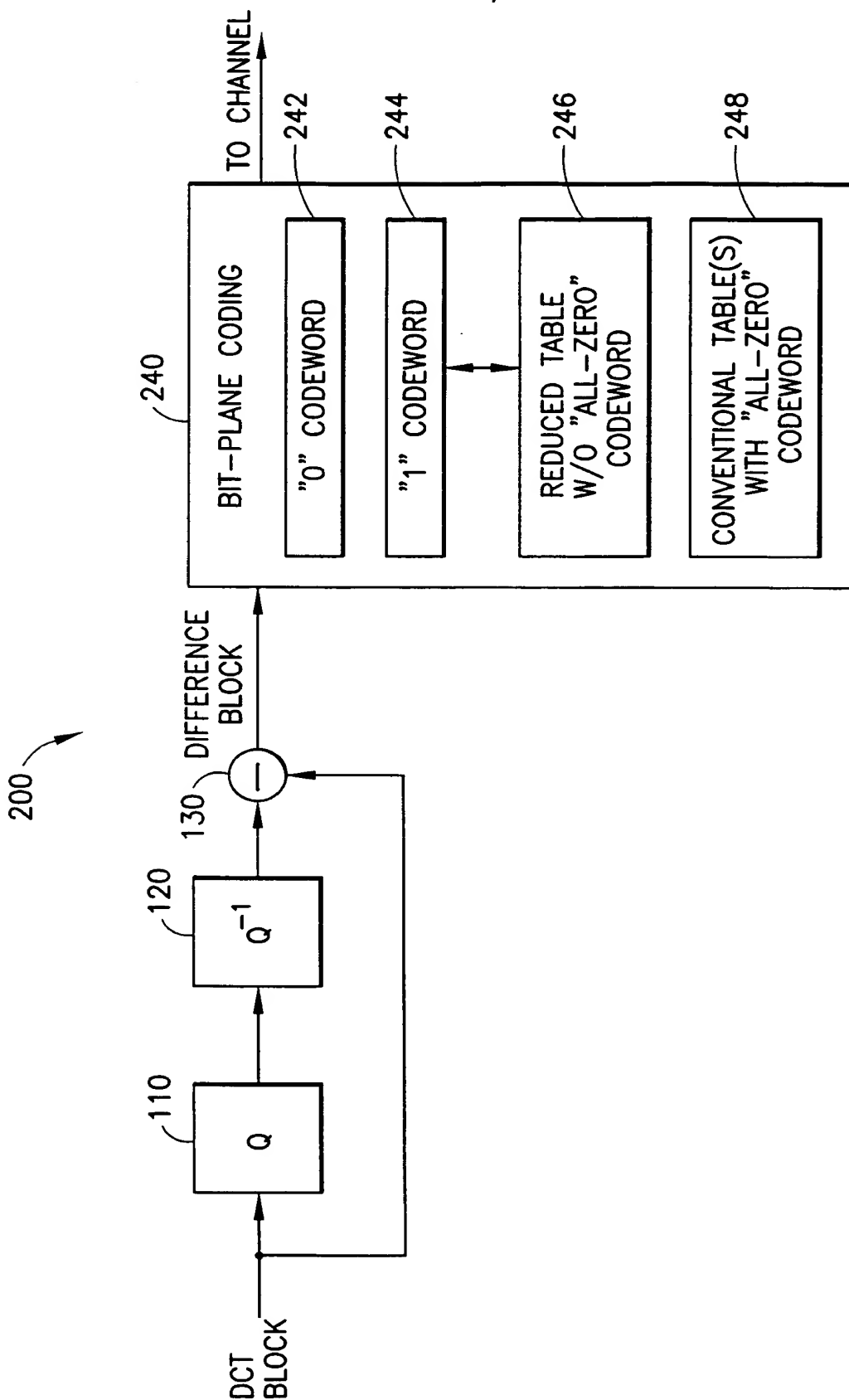


FIG.2

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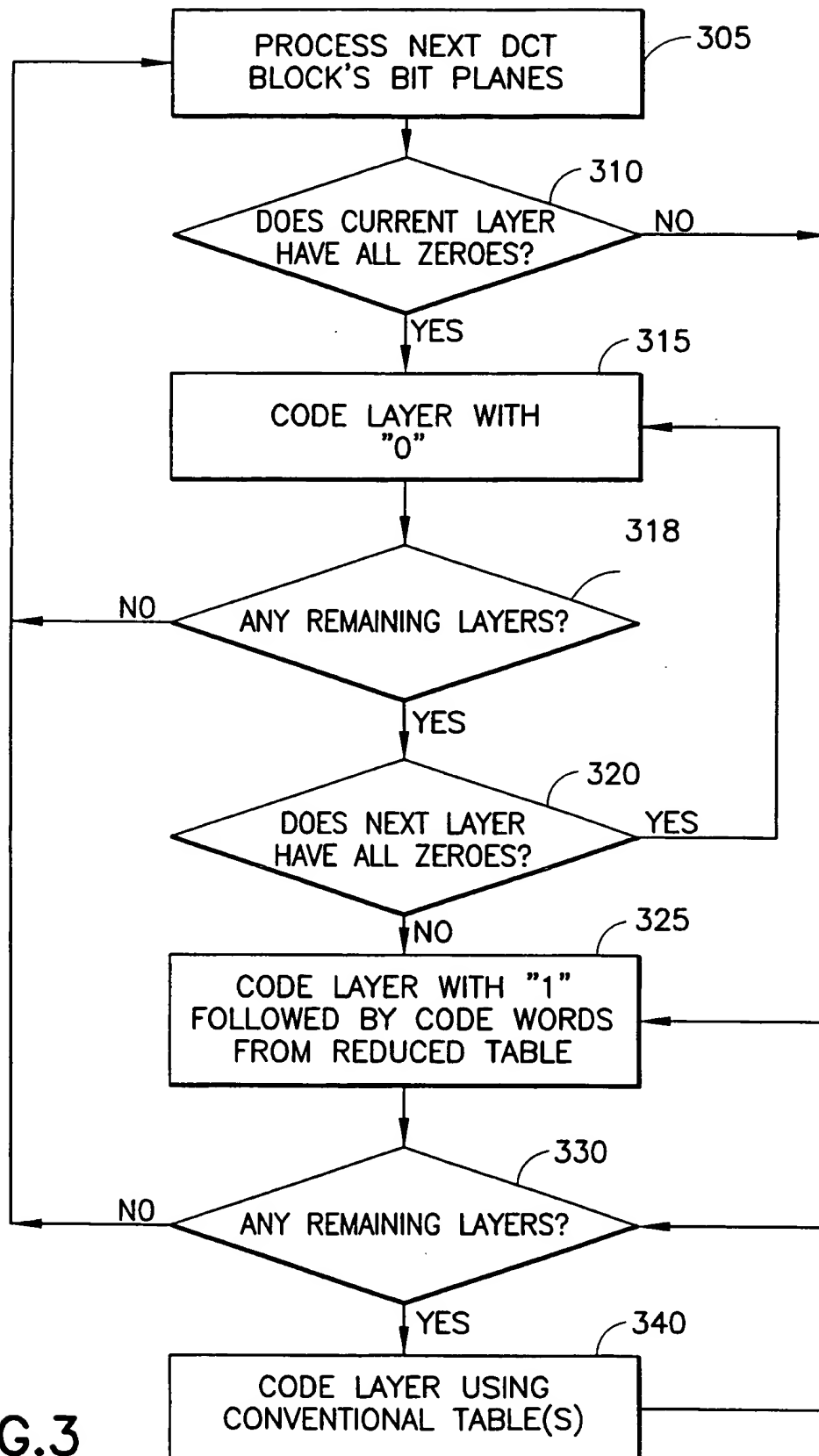


FIG. 3

4/4

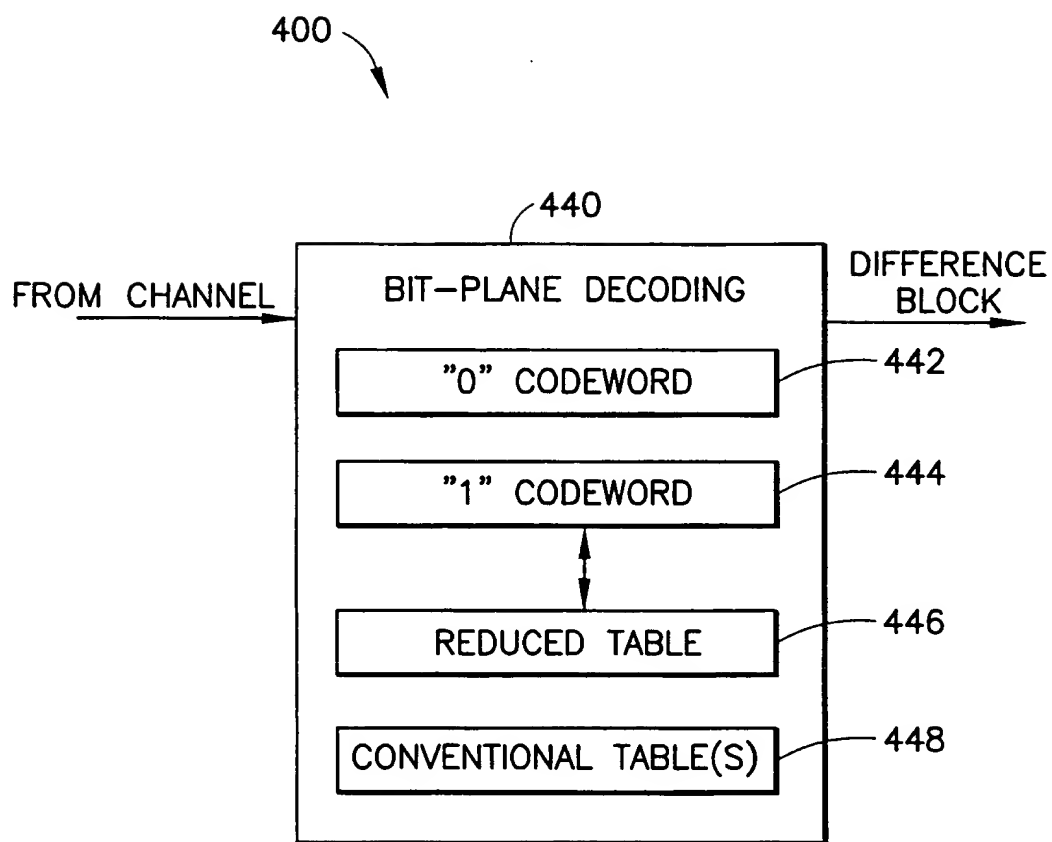


FIG.4

INTERNATIONAL SEARCH REPORT

International Application No
PCT/US 99/25973

A. CLASSIFICATION OF SUBJECT MATTER
IPC 7 H04N1/41 H04N7/26

According to International Patent Classification (IPC) or to both national classification and IPC

B. FIELDS SEARCHED

Minimum documentation searched (classification system followed by classification symbols)
IPC 7 H04N

Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched

Electronic data base consulted during the international search (name of data base and, where practical, search terms used)

C. DOCUMENTS CONSIDERED TO BE RELEVANT

Category *	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
X	EP 0 855 838 A (CANON INFORMATION SYSTEMS RESEARCH AUSTRALIA PTY LTD.) 29 July 1998 (1998-07-29) page 12, line 15 -page 15, line 44; figures 1,3-6	1,2,6-8, 12-16, 19,20
A,P	FAN LING; WEIPING LI; HONGQIAO SUN : "Bitplane coding of DCT coefficients for image and video compression " PROCEEDINGS OF THE SPIE, vol. 3653, 27 January 1999 (1999-01-27), pages 500-508, XP002133777 page 503, paragraph 2 -page 505, paragraph 1 --- -/-	1-22

☒ Further documents are listed in the continuation of box C.

☒ Patent family members are listed in annex.

* Special categories of cited documents :

- *A* document defining the general state of the art which is not considered to be of particular relevance
- *E* earlier document but published on or after the international filing date
- *L* document which may throw doubts on priority claim(s) or which is cited to establish the publication date of another citation or other special reason (as specified)
- *O* document referring to an oral disclosure, use, exhibition or other means
- *P* document published prior to the international filing date but later than the priority date claimed

- *T* later document published after the international filing date or priority date and not in conflict with the application but cited to understand the principle or theory underlying the invention
- *X* document of particular relevance; the claimed invention cannot be considered novel or cannot be considered to involve an inventive step when the document is taken alone
- *Y* document of particular relevance; the claimed invention cannot be considered to involve an inventive step when the document is combined with one or more other such documents, such combination being obvious to a person skilled in the art.
- *&* document member of the same patent family

Date of the actual completion of the international search

28 March 2000

Date of mailing of the international search report

14/04/2000

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Authorized officer

Raeymaekers, P

INTERNATIONAL SEARCH REPORT

International Application No
PCT/US 99/25973

C.(Continuation) DOCUMENTS CONSIDERED TO BE RELEVANT

Category *	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
A	<p>FAN LING, WEIPING LI: "dimensional adaptive arithmetic coding for image compression"</p> <p>PROCEEDINGS OF THE 1998 IEEE INTERNATIONAL SYMPOSIUM ON CIRCUITS AND SYSTEMS, 'Online!</p> <p>vol. 4, 3 June 1998 (1998-06-03), pages 13-16, XP002133776</p> <p>Retrieved from the Internet: <URL:http://iel.ihs.com> 'retrieved on 2000-03-22!</p> <p style="text-align: center;">-----</p>	1-22

INTERNATIONAL SEARCH REPORT

Information on patent family members

International Application No

PCT/US 99/25973

Patent document cited in search report	Publication date	Patent family member(s)	Publication date
EP 855838 A	29-07-1998	AU 714202 B	23-12-1999
		AU 5270298 A	30-07-1998
		JP 10313457 A	24-11-1998
<hr/>			

PCT

REC'D 16 JAN 2001

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PCT

INTERNATIONAL PRELIMINARY EXAMINATION REPORT

(PCT Article 36 and Rule 70)

Applicant's or agent's file reference GIC-562 PCT	FOR FURTHER ACTION See Notification of Transmittal of International Preliminary Examination Report (Form PCT/IPEA/416)	
International application No. PCT/US99/25973	International filing date (day/month/year) 04 NOVEMBER 1999	Priority date (day/month/year) 04 DECEMBER 1998
International Patent Classification (IPC) or national classification and IPC IPC(7): H04N 1/41, 7/26 and US Cl.: 375/240.19		
Applicant GENERAL INSTRUMENT CORPORATION		

1. This international preliminary examination report has been prepared by this International Preliminary Examining Authority and is transmitted to the applicant according to Article 36.
2. This REPORT consists of a total of 4 sheets.
- ☒ This report is also accompanied by ANNEXES, i.e., sheets of the description, claims and/or drawings which have been amended and are the basis for this report and/or sheets containing rectifications made before this Authority. (see Rule 70.16 and Section 607 of the Administrative Instructions under the PCT).
- These annexes consist of a total of 4 sheets.

3. This report contains indications relating to the following items:

- I ☒ Basis of the report
- II ☐ Priority
- III ☐ Non-establishment of report with regard to novelty, inventive step or industrial applicability
- IV ☐ Lack of unity of invention
- V ☒ Reasoned statement under Article 35(2) with regard to novelty, inventive step or industrial applicability; citations and explanations supporting such statement
- VI ☐ Certain documents cited
- VII ☐ Certain defects in the international application
- VIII ☐ Certain observations on the international application

Date of submission of the demand 06 JUNE 2000	Date of completion of this report 14 DECEMBER 2000
Name and mailing address of the IPEA/US Commissioner of Patents and Trademarks Box PCT Washington, D.C. 20231 Facsimile No. (703) 305-3230	Authorized officer CHRISTOPHER KELLY Telephone No. (703) 305-3230

INTERNATIONAL PRELIMINARY EXAMINATION REPORT

International application No.

PCT/US99/25973

I. Basis of the report

1. With regard to the elements of the international application: *

☐ the international application as originally filed☒ the description:

pages (See Attached) , as originally filed
pages , filed with the demand
pages , filed with the letter of

☒ the claims:

pages (See Attached) , as originally filed
pages , as amended (together with any statement) under Article 19
pages , filed with the demand
pages , filed with the letter of

☒ the drawings:

pages (See Attached) , as originally filed
pages , filed with the demand
pages , filed with the letter of

☒ the sequence listing part of the description:

pages (See Attached) , as originally filed
pages , filed with the demand
pages , filed with the letter of

2. With regard to the language, all the elements marked above were available or furnished to this Authority in the language in which the international application was filed, unless otherwise indicated under this item.

These elements were available or furnished to this Authority in the following language _____ which is:

- ☐ the language of a translation furnished for the purposes of international search (under Rule 23.1(b)).
☐ the language of publication of the international application (under Rule 48.3(b)).
☐ the language of the translation furnished for the purposes of international preliminary examination (under Rules 55.2 and/or 55.3).

3. With regard to any nucleotide and/or amino acid sequence disclosed in the international application, the international preliminary examination was carried out on the basis of the sequence listing:

- ☐ contained in the international application in printed form.
☐ filed together with the international application in computer readable form.
☐ furnished subsequently to this Authority in written form.
☐ furnished subsequently to this Authority in computer readable form.
☐ The statement that the subsequently furnished written sequence listing does not go beyond the disclosure in the international application as filed has been furnished.
☐ The statement that the information recorded in computer readable form is identical to the written sequence listing has been furnished.

4. ☒ The amendments have resulted in the cancellation of:

- ☒ the description, pages NONE
☒ the claims, Nos. NONE
☒ the drawings, sheets/fig NONE

5. ☐ This report has been drawn as if (some of) the amendments had not been made, since they have been considered to go beyond the disclosure as filed, as indicated in the Supplemental Box (Rule 70.2(c)).**

* Replacement sheets which have been furnished to the receiving Office in response to an invitation under Article 14 are referred to in this report as "originally filed" and are not annexed to this report since they do not contain amendments (Rules 70.16 and 70.17).

**Any replacement sheet containing such amendments must be referred to under item 1 and annexed to this report.

V. Reasoned statement under Article 35(2) with regard to novelty, inventive step or industrial applicability; citations and explanations supporting such statement**1. statement**

Novelty (N)	Claims <u>3-4, 9-10, 17-18, 21-22</u>	YES
	Claims <u>1-2, 5-8, 11-16, 19-20</u>	NO
Inventive Step (IS)	Claims <u>3-4, 9-10, 17-18, 21-22</u>	YES
	Claims <u>1-2, 5-8, 11-16, 19-20</u>	NO
Industrial Applicability (IA)	Claims <u>1-22</u>	YES
	Claims <u>NONE</u>	NO

2. citations and explanations (Rule 70.7)

Claims 1-2, 5-8, 11-16, and 19-20 lack novelty under PCT Article 33(2) as being anticipated by CANON INFORMATION SYSTEMS RESEARCH AUSTRALIA PTY LTD.

Regarding claims 1-2, 7-8, 13-16, and 19-20, CANON INFORMATION SYSTEMS RESEARCH AUSTRALIA PTY LTD teaches the same apparatus and method for efficient coding of a plurality of bit planes in which transform coefficient data is carried (See page 2, lines 32-37), comprising the steps of providing a codeword having a first state designating that a bit plane comprises all binary zeroes (See page 2, lines 55-58, and page 3, lines 1-2), and a second state designating that a bit plane does not comprise all binary zeroes (See page 3, lines 1-2), coding the most significant bit (MSB) bit plane with the codeword in its first state when the MSB bit plane comprises all zeroes (See page 12, lines 22-33), proceeding from the MSB bit plane toward the least significant bit (LSB) bit plane, coding each successive bit plane that comprises all binary zeroes with the codeword in its first state, if any such successive bit planes are present (See page 9, lines 42-50), until a first bit plane that does not comprise all binary zeroes is reached (See page 8, lines 2-5), and coding the first bit plane with the codeword in its second state, followed by at least one codeword obtained from a first entropy coding table according to the bits in the first bit plane (See page 14, lines 46-58), and wherein the codeword having the first and second states is a one-bit codeword (See page 12, line 54-56).

As per claims 5-6, and 11-12, CANON INFORMATION SYSTEMS RESEARCH AUSTRALIA PTY LTD further teaches the same apparatus and method for efficient coding of a plurality of bit planes wherein the transform coefficients comprised Discrete Cosine transform (DCT) data (See page 11, lines 1-3), and wherein the transform coefficients comprises image data (See page 15, line 51-52, and fig. 9).

(Continued on Supplemental Sheet.)

Supplemental Box

(To be used when the space in any of the preceding boxes is not sufficient)

Continuation of: Boxes I - VIII

Sheet 10

I. BASIS OF REPORT:

This report has been drawn on the basis of the description,
page(s) 1-21, as originally filed.
page(s) NONE, filed with the demand.
and additional amendments:
NONE

This report has been drawn on the basis of the claims,
page(s) 23, 25, and 27 as originally filed.
page(s) NONE, as amended under Article 19.
page(s) NONE, filed with the demand.
and additional amendments:
Pages 22, 24, 26, and 28, filed with the letter of 02 August 2000.

This report has been drawn on the basis of the drawings,
page(s) 1-4, as originally filed.
page(s) NONE, filed with the demand.
and additional amendments:
NONE

This report has been drawn on the basis of the sequence listing part of the description:
page(s) NONE, as originally filed.
pages(s) NONE, filed with the demand.
and additional amendments:
NONE

V. 2. REASONED STATEMENTS - CITATIONS AND EXPLANATIONS (Continued):

Claims 3-4, 9-10, 17-18, and 21-22 meet the criteria set out in PCT Article 33(2)-(4), because the prior art does not teach or fairly suggest an apparatus and method for efficient coding of a plurality of bit planes in which transform coefficient data is carried, comprising the steps of providing a codeword having a first state designating that a bit plane comprises all binary zeroes, and a second state designating that a bit plane does not comprise all binary zeroes, coding the most significant bit (MSB) bit plane with the codeword in its first state when the MSB bit plane comprises all zeroes, proceeding from the MSB bit plane toward the least significant bit (LSB) bit plane, coding each successive bit plane that comprises all binary zeroes with the codeword in its first state, if any such successive bit planes are present, until a first bit plane that does not comprise all binary zeroes is reached, and coding the first bit plane with the codeword in its second state, followed by at least one codeword obtained from a first entropy coding table according to the bits in the first bit plane, and wherein the first entropy coding table does not include a multi-bit codeword for coding a bit plane with all binary zeroes, and a second entropy coding table which includes a multibit codeword for coding a bit plane with all binary zeroes.

----- NEW CITATIONS -----

US 5,029,122 A (UETANI) 02 JULY 1991, see column 16, lines 30-60.

US 5,768,535 A (CHADDHA et al) 16 JUNE 1998, see column 10, lines 47-53.

What is claimed is:

1. A method for efficient coding of a plurality of bit planes in which transform coefficient data is carried, comprising the steps of:

(a) providing a codeword having a first state designating that a bit plane comprises all binary zeroes, and a second state designating that a bit plane does not comprise all binary zeroes;

(b) coding the most significant bit (MSB) bit plane with said codeword in its first state when the MSB bit plane comprises all binary zeroes;

(c) proceeding from the MSB bit plane toward the least significant bit (LSB) bit plane, coding each successive bit plane that comprises all binary zeroes with said codeword in its first state, if any such successive bit planes are present, until a first bit plane that does not comprise all binary zeroes is reached; and

(d) coding said first bit plane with said codeword in its second state, followed by at least one codeword obtained from a first entropy coding table according to the bits in said first bit plane.

2. The method of claim 1, wherein:

the codeword having the first and second states is a one-bit codeword.

3. The method of claim 1, wherein:

the first entropy coding table does not include a multi-bit codeword for coding a bit plane with all binary zeroes.

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codeword in its first state for each successive bit plane that comprises all binary zeroes, if any such successive bit planes are present, until a first bit plane that does not comprise all binary zeroes is reached; and

(c) decoding said codeword in its second state for said first bit plane, then using a first entropy decoding table for decoding at least one codeword that follows said codeword in its second state;

wherein the at least one codeword is obtained from a first entropy coding table according to the bits in said first bit plane.

8. The method of claim 7, wherein:
the codeword having the first and second states is a one-bit codeword.

9. The method of claim 7, wherein:
the first entropy decoding table does not include a multi-bit codeword for decoding a bit plane with all binary zeroes.

10. The method of claim 7, comprising the further step of:

providing a second entropy decoding table for decoding at least one bit plane that follows the first bit plane; wherein:

the second entropy decoding table includes a multi-bit codeword for decoding a bit plane with all binary zeroes.

coding table according to the bits in said first bit plane.

14. The signal of claim 13, wherein:
the codeword having the first and second states is a one-bit codeword.

15. An apparatus for efficient coding of a plurality of bit planes in which transform coefficient data is carried, comprising:

(a) means for providing a codeword having a first state designating that a bit plane comprises all binary zeroes, and a second state designating that a bit plane does not comprise all binary zeroes;

(b) means for coding the most significant bit (MSB) bit plane with said codeword in its first state when the MSB bit plane comprises all binary zeroes;

(c) means for coding each successive bit plane that comprises all binary zeroes with said codeword in its first state, if any such successive bit planes are present, proceeding from the MSB bit plane toward the least significant bit (LSB) bit plane, until a first bit plane that does not comprise all binary zeroes is reached; and

(d) means for coding said first bit plane with said codeword in its second state, followed by at least one codeword obtained from a first entropy coding table according to the bits in said first bit plane.

16. The apparatus of claim 15, wherein:
the codeword having the first and second states is a one-bit codeword.

(c) means for decoding said codeword in its second state for said first bit plane, then using a first entropy decoding table for decoding at least one codeword that follows said codeword in its second state;

wherein the at least one codeword is obtained from a first entropy coding table according to the bits in said first bit plane.

20. The apparatus of claim 19, wherein:
the codeword having the first and second states is a one-bit codeword.

21. The apparatus of claim 19, wherein:
the first entropy decoding table does not include a multi-bit codeword for decoding a bit plane with all binary zeroes.

22. The apparatus of claim 19, further comprising:

a second entropy decoding table for decoding at least one bit plane that follows the first bit plane;
wherein:

the second entropy decoding table includes a multi-bit codeword for decoding a bit plane with all binary zeroes.

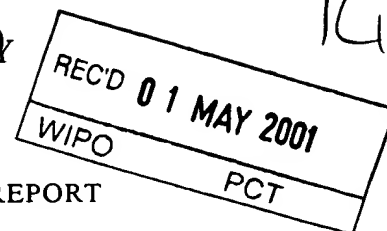
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PATENT COOPERATION TREATY

PCT

INTERNATIONAL PRELIMINARY EXAMINATION REPORT

(PCT Article 36 and Rule 70)



Applicant's or agent's file reference GIC-562 PCT	FOR FURTHER ACTION See Notification of Transmittal of International Preliminary Examination Report (Form PCT/IPEA/416)	
International application No. PCT/US99/25973	International filing date (day/month/year) 04 NOVEMBER 1999	Priority date (day/month/year) 04 DECEMBER 1998
International Patent Classification (IPC) or national classification and IPC IPC(7): H04N 1/41, 7/26 and US Cl.: 375/240.19		
Applicant GENERAL INSTRUMENT CORPORATION		

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- This international preliminary examination report has been prepared by this International Preliminary Examining Authority and is transmitted to the applicant according to Article 36.
- This REPORT consists of a total of 4 sheets.
☒ This report is also accompanied by ANNEXES, i.e., sheets of the description, claims and/or drawings which have been amended and are the basis for this report and/or sheets containing rectifications made before this Authority. (see Rule 70.16 and Section 607 of the Administrative Instructions under the PCT).
 These annexes consist of a total of 4 sheets.

3. This report contains indications relating to the following items:

- I ☒ Basis of the report
- II ☐ Priority
- III ☐ Non-establishment of report with regard to novelty, inventive step or industrial applicability
- IV ☐ Lack of unity of invention
- V ☒ Reasoned statement under Article 35(2) with regard to novelty, inventive step or industrial applicability; citations and explanations supporting such statement
- VI ☐ Certain documents cited
- VII ☐ Certain defects in the international application
- VIII ☐ Certain observations on the international application

**CORRECTED
VERSION**

Date of submission of the demand 06 JUNE 2000	Date of completion of this report 03 APRIL 2001
Name and mailing address of the IPEA/US Commissioner of Patents and Trademarks Box PCT Washington, D.C. 20231	Authorized officer CHRISTOPHER KELLY <i>Kygenia Zogor</i>
Facsimile No. (703) 305-3230	Telephone No. (703) 305-4700

INTERNATIONAL PRELIMINARY EXAMINATION REPORT

International application No.

PCT/US99/25973

I. Basis of the report**1. With regard to the elements of the international application:***

- ☐ the international application as originally filed
- ☒ the description:
pages _____ (See Attached) _____, as originally filed
pages _____, filed with the demand
pages _____, filed with the letter of _____
- ☒ the claims:
pages _____ (See Attached) _____, as originally filed
pages _____, as amended (together with any statement) under Article 19
pages _____, filed with the demand
pages _____, filed with the letter of _____
- ☒ the drawings:
pages _____ (See Attached) _____, as originally filed
pages _____, filed with the demand
pages _____, filed with the letter of _____
- ☒ the sequence listing part of the description:
pages _____ (See Attached) _____, as originally filed
pages _____, filed with the demand
pages _____, filed with the letter of _____

2. With regard to the language, all the elements marked above were available or furnished to this Authority in the language in which the international application was filed, unless otherwise indicated under this item.

These elements were available or furnished to this Authority in the following language _____ which is:

- ☐ the language of a translation furnished for the purposes of international search (under Rule 23.1(b)).
- ☐ the language of publication of the international application (under Rule 48.3(b)).
- ☐ the language of the translation furnished for the purposes of international preliminary examination (under Rules 55.2 and/or 55.3).

3. With regard to any nucleotide and/or amino acid sequence disclosed in the international application, the international preliminary examination was carried out on the basis of the sequence listing:

- ☐ contained in the international application in printed form.
- ☐ filed together with the international application in computer readable form.
- ☐ furnished subsequently to this Authority in written form.
- ☐ furnished subsequently to this Authority in computer readable form.
- ☐ The statement that the subsequently furnished written sequence listing does not go beyond the disclosure in the international application as filed has been furnished.
- ☐ The statement that the information recorded in computer readable form is identical to the written sequence listing has been furnished.

4. ☒ The amendments have resulted in the cancellation of:

- ☒ the description, pages _____ none _____
- ☒ the claims, Nos. _____ NONE _____
- ☒ the drawings, sheets/fig _____ NONE _____

5. ☐ This report has been drawn as if (some of) the amendments had not been made, since they have been considered to go beyond the disclosure as filed, as indicated in the Supplemental Box (Rule 70.2(c)).**

* Replacement sheets which have been furnished to the receiving Office in response to an invitation under Article 14 are referred to in this report as "originally filed" and are not annexed to this report since they do not contain amendments (Rules 70.16 and 70.17).

**Any replacement sheet containing such amendments must be referred to under item 1 and annexed to this report

INTERNATIONAL PRELIMINARY EXAMINATION REPORT

International application No.

PCT/US99/25973

V. Reasoned statement under Article 35(2) with regard to novelty, inventive step or industrial applicability; citations and explanations supporting such statement**1. statement**

Novelty (N)	Claims <u>1-22</u>	YES
	Claims <u>NONE</u>	NO
Inventive Step (IS)	Claims <u>1-22</u>	YES
	Claims <u>NONE</u>	NO
Industrial Applicability (IA)	Claims <u>1-22</u>	YES
	Claims <u>NONE</u>	NO

2. citations and explanations (Rule 70.7)

Claims 1-22 meet the criteria set out in PCT Article 33(2)-(4), because the prior art does not teach or fairly suggest an apparatus and method for efficient coding of a plurality of bit planes in which transform coefficient data is carried, comprising the steps of providing a codeword having a first state designating that a bit plane comprises all binary zeroes, and a second state designating that a bit plane does not comprise all binary zeroes, coding the most significant bit (MSB) bit plane with the codeword in its first state when the MSB bit plane comprises all zeroes, proceeding from the MSB bit plane toward the least significant bit (LSB) bit plane, coding each successive bit plane that comprises all binary zeroes with the codeword in its first state, if any such successive bit planes are present, until a first bit plane that does not comprise all binary zeroes is reached, and coding the first bit plane with the codeword in its second state, followed by at least one codeword obtained from a first entropy coding table according to the bits in the first bit plane, and wherein the first entropy coding table does not include a multi-bit codeword for coding a bit plane with all binary zeroes, and a second entropy coding table which includes a multibit codeword for coding a bit plane with all binary zeroes.

----- NEW CITATIONS -----

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US 5,768,535 A (CHADDHA et al) 16 JUNE 1998, see column 10, lines 47-53.

INTERNATIONAL PRELIMINARY EXAMINATION REPORT

International application No.

PCT/US99/25973

Supplemental Box

(To be used when the space in any of the preceding boxes is not sufficient)

Continuation of: Boxes I - VIII

Sheet 10

I. BASIS OF REPORT:

This report has been drawn on the basis of the description,
page(s) 1-21, as originally filed.
page(s) NONE, filed with the demand.
and additional amendments:
NONE

This report has been drawn on the basis of the claims,
page(s) 23, 25 and 27, as originally filed.
page(s) NONE, as amended under Article 19.
page(s) NONE, filed with the demand.
and additional amendments:
Pages 22, 24, 26, and 28, filed with the letter of 02 August 2000.

This report has been drawn on the basis of the drawings,
page(s) 1-4, as originally filed.
page(s) NONE, filed with the demand.
and additional amendments:
NONE

This report has been drawn on the basis of the sequence listing part of the description:
page(s) NONE, as originally filed.
pages(s) NONE, filed with the demand.
and additional amendments:
NONE

What is claimed is:

1. A method for efficient coding of a plurality of bit planes in which transform coefficient data is carried, comprising the steps of:

(a) providing a codeword having a first state designating that a bit plane comprises all binary zeroes, and a second state designating that a bit plane does not comprise all binary zeroes;

(b) coding the most significant bit (MSB) bit plane with said codeword in its first state when the MSB bit plane comprises all binary zeroes;

(c) proceeding from the MSB bit plane toward the least significant bit (LSB) bit plane, coding each successive bit plane that comprises all binary zeroes with said codeword in its first state, if any such successive bit planes are present, until a first bit plane that does not comprise all binary zeroes is reached; and

(d) coding said first bit plane with said codeword in its second state, followed by at least one codeword obtained from a first entropy coding table according to the bits in said first bit plane.

2. The method of claim 1, wherein:
the codeword having the first and second states is a one-bit codeword.

3. The method of claim 1, wherein:
the first entropy coding table does not include a multi-bit codeword for coding a bit plane with all binary zeroes.

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codeword in its first state for each successive bit plane that comprises all binary zeroes, if any such successive bit planes are present, until a first bit plane that does not comprise all binary zeroes is reached; and

(c) decoding said codeword in its second state for said first bit plane, then using a first entropy decoding table for decoding at least one codeword that follows said codeword in its second state;

wherein the at least one codeword is obtained from a first entropy coding table according to the bits in said first bit plane.

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9. The method of claim 7, wherein:
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coding table according to the bits in said first bit plane.

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22. The apparatus of claim 19, further comprising:

a second entropy decoding table for decoding at least one bit plane that follows the first bit plane;
wherein:

the second entropy decoding table includes a multi-bit codeword for decoding a bit plane with all binary zeroes.

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